

Conflict-Free Coloring of Extended Duplicate Graph of Twig

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Abstract. Graph coloring is one of the most important area of research in graph theory. Conflict-free coloring is defined as a vertex coloring such that for each vertex $v \in V$, there exists a vertex in the neighbourhood of v denoted by $N(v)$, whose color is different from color of each other vertex in $N[v]$. In this paper, we prove the existence of conflict-free vertex coloring in $EDG(T_m)$ and show that $\chi_{CF}(EDG(T_m)) = \Delta + 1$. As a consequence, we show that the induced subgraph of each color class with s vertices is $\overline{K_s}$, where $\overline{K_s}$ is the complement of the complete graph K_s .

Keywords: Conflict-free coloring, extended duplicate graph of twig graphs.

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1. Introduction

The vertex coloring of a graph is coloring the vertices of the graph in such a way that adjacent vertices have different colors. Motivated by the problem of frequency assignments in cellular networks, Even et.al [3] and Smorodinsky [6] introduced the concept of conflict-free coloring. Pach and Tarados [5] instituted the idea of conflict coloring through graphs and hyper graphs. Glebov et.al [4] further brought forward the concept of conflict-free coloring through simple graphs. The Extended Duplicate Graph of Twig graphs was introduced by Thirusangu et.al [7]. Other kind of labelling are studied in [8-13].

In this paper, we prove the existence of conflict-free vertex coloring in Extended Duplicate Graph of Twig graphs $EDG(T_m)$ which yields the result that the induced subgraph of each color class with s vertices is $\overline{K_s}$, where $\overline{K_s}$ is the complement of the complete graph K_s . It is also proved that $\chi_{CF}(EDG(T_m)) = \Delta + 1$.

2. Preliminaries

In this section we give the basic notions relevant to this paper.

Definition 2.1. A graph $G(V,E)$ obtained from a path by joining exactly two pendent edges to each internal vertices of a path is called a Twig graph. A Twig graph with m internal vertices is denoted by T_m .

Definition 2.2. Let $G(V,E)$ be a simple graph. A duplicate graph of G is $DG=(V_1, E_1)$ where the vertex set $V_1 = V \cup V'$ and $V \cap V' = \emptyset$ and $f:V \rightarrow V'$ is bijective and the edgeset E_1 of DG is defined as follows: The edge uv is in E if and only if both uv and $u'v'$ are edges in E_1 .

Definition 2.3. Let $DG=(V_1, E_1)$ be a duplicate graph of the Twig graph $G(V,E)$. We add an edge between any one vertex from V to any other vertex in V' except the terminal vertices of V and V' . For our convenience, we take $v_2 \in V$ and $v_2' \in V'$ and thus the edge v_2v_2' is formed. We call this new graph as the extended duplicate graph of the Twig graph T_m and is denoted by $EDG(T_m)$. The vertex set and the edge set of $EDG(T_m)$ are given as follows :

$$V = \{v_1, v_2, \dots, v_{3m+2}, v_1', v_2', \dots, v_{3m+2}'\}$$

$$E = \{v_1v_2', v_1'v_2, v_2v_2'\} \cup \{v_s'v_l \cup v_{s+3}v_{l+3}' \cup v_s v_l' \cup v_{s+3}v_{l+3}' / s=2+6i, l = s+j+1; 0 \leq i \leq \lfloor \frac{m-2}{2} \rfloor; 0 \leq j \leq 2\}$$

$$E = \{v_1v_2', v_1'v_2, v_2v_2'\} \cup \{v_s'v_l \cup v_s v_l' / s=2+6i, l = s+j+1; 0 \leq i \leq \lfloor \frac{m-1}{2} \rfloor; 0 \leq j \leq 2\} \cup \{v_{s+3}'v_{l+3} \cup v_{s+3}v_{l+3}' / s=2+6i, l = s+j+1; 0 \leq i \leq \lfloor \frac{m-3}{2} \rfloor; 0 \leq j \leq 2\}$$

Clearly $EDG(T_m)$ has $6m+4$ vertices and $6m+3$ edges where m is the number of internal vertices of the Twig graph.

Definition 2.4. Let $G=(V,E)$ be a simple graph. For every vertex $u \in V$, we denote $N(u) = \{v \in V: uv \in E\}$, its neighbourhood and by $N(u) \cup \{u\} = N[u]$, its closed neighbourhood. Also the maximum degree of the graph is denoted by Δ .

Definition 2.5. A vertex coloring χ of G is called conflict-free if for each vertex $u \in V$, there exists a vertex v in $N(u)$ whose color is different from the color of each other vertex in $N[u]$.

The conflict-free chromatic number $\chi_{CF}(G)$ is the smallest r , such that there exists a conflict-free r -coloring of G .

Definition 2.6. For any set S of vertices of graph G , the induced subgraph $\langle S \rangle$, is the maximal subgraph of G with vertex set S . Thus two vertices are adjacent in S if and only if they are adjacent in G .

Definition 2.7. A graph G is said to be a complete graph if each pair of vertices is adjacent. A complete graph of n vertices is denoted by K_n .

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The complement \bar{G} of G is defined to be the graph which has V as its vertex set and two vertices are adjacent in \bar{G} if and only if they are not adjacent in G .

3. Main results

In this section, we present an algorithm to obtain the conflict-free coloring of $EDG(T_m)$. We prove that $EDG(T_m)$ admits conflict-free coloring. We also show that the induced subgraph of each of the color class is a totally disconnected graph. Also we find the conflict-free chromatic number of $EDG(T_m)$ is $\Delta + 1$.

Algorithm 3.1:

Procedure: Conflict-free coloring of $EDG(T_m)$

Input: $V \leftarrow \{v_1, v_2, \dots, v_{3m+2}, v_1', v_2', \dots, v_{3m+2}'\}$

$E \leftarrow \{e_1, e_2, \dots, e_{3m+2}, e_1', e_2', \dots, e_{3m+1}'\}$

Output: Conflict-free colored $EDG(T_m)$ graph.

if $m=1$

$$v_1 \cup v_1' \leftarrow C_1$$

$$v_2 \leftarrow C_2$$

$$v_2' \leftarrow C_3$$

$$v_3 \cup v_3' \leftarrow C_4$$

else

if $m=2$

$$v_1 \cup v_1' \cup v_6 \cup v_6' \leftarrow C_1$$

$$v_2 \cup v_7' \leftarrow C_2$$

$$v_2' \cup v_7 \leftarrow C_3$$

$$v_3 \cup v_3' \cup v_8 \cup v_8' \leftarrow C_4$$

else

for $n= 2$ to m do

$$v_1 \cup v_1' \cup (\cup_{i=2}^n v_{3i} \cup v_{3i}') \leftarrow C_1$$

$$v_2 \cup v_7' \cup (\cup_{i=3}^n v_{3i+1} \cup v_{3i+1}') \leftarrow C_2$$

$$v_2' \cup v_7 \cup (\cup_{i=3k/k \in N}^n v_{3i+2} \cup v_{3i+2}') \leftarrow C_3$$

$$v_3 \cup v_3' \cup (\cup_{i=3k-1/k \in N}^n v_{3i+2} \cup v_{3i+2}') \leftarrow C_4$$

end for

end if

end if

if $(m \leq 3)$

$$v_4 \cup v_4' \leftarrow C_5$$

else

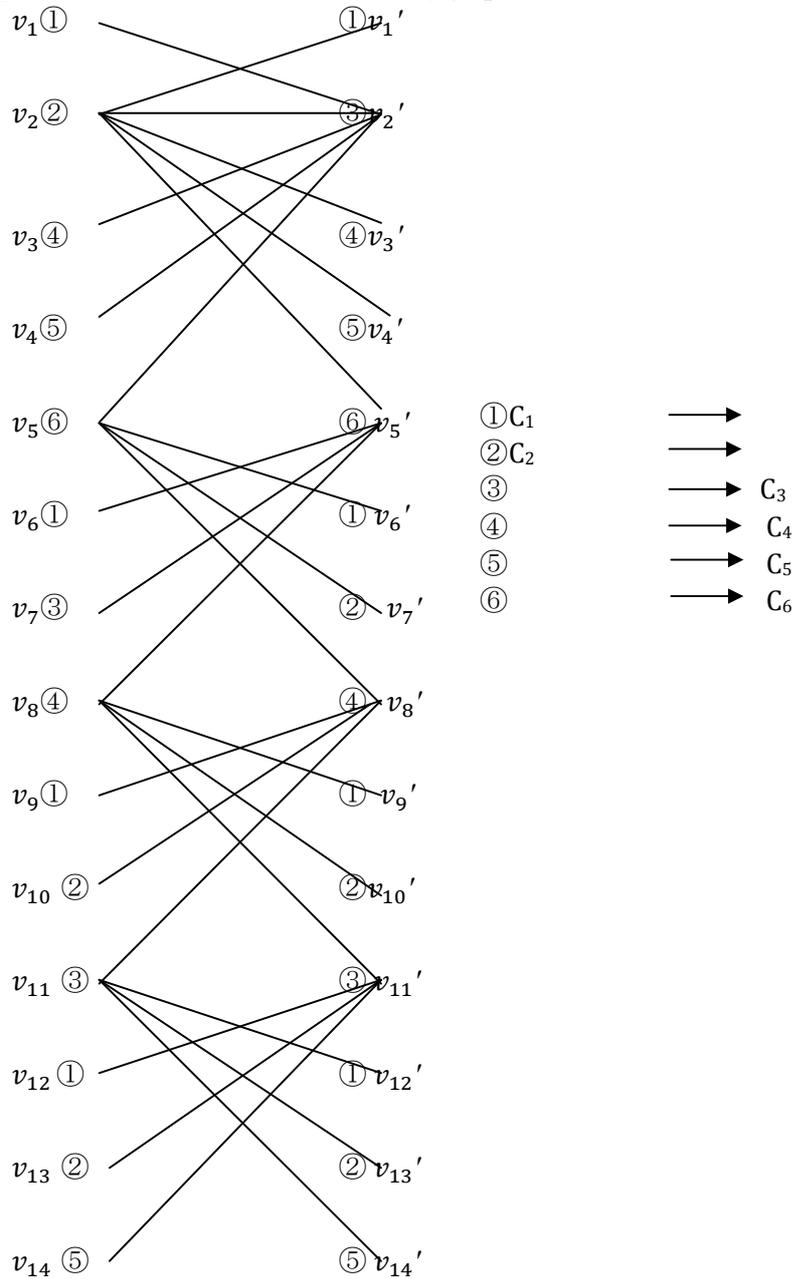
$$v_4 \cup v_4' \cup (\cup_{i=3k+1/k \in N}^m v_{3i+2} \cup v_{3i+2}') \leftarrow C_5$$

end if

$$v_5 \cup v_5' \leftarrow C_6$$

end procedure

Example 3.1. Conflict-free colored EDG(T_4) graph.



Here C_i for each i represent the color, $1 \leq i \leq 6$.

Figure 1:

Conflict-Free Coloring of Extended Duplicate Graph of Twig

Theorem 3.1. $EDG(T_m)$ is a conflict-free colorable graph.

Proof: The $EDG(T_m)$ has $6m+4$ vertices and $6m+3$ edges.

To prove $EDG(T_m)$ is a conflict-free colorable graph, it is enough to prove that the vertex of maximum degree, either v_2 or v_2' is conflict-free colorable.

Using Algorithm 3.1, the color class C_2 has vertices

$v_2, v_7, v_{10}, v_{10}, v_{13}, v_{13}, \dots, v_{3m+1}, v_{3m+1}$. Thus $v_2 \in C_2$. Also the neighbours of v_2 are $\{v_1', v_2', v_3', v_4', v_5'\}$ where $v_1' \in C_1, v_2' \in C_3, v_3' \in C_4, v_4' \in C_5, v_5' \in C_6$. That is each neighbours of v_2 is in different color classes.

Similarly for all the vertices of $EDG(T_m)$ in a color class has neighbours in different color classes.

Hence $EDG(T_m)$ is a conflict-free colorable graph.

Theorem 3.2. In Extended Duplicate Graph $EDG(T_m)$, the induced subgraph of each color class with n_i vertices is totally disconnected graph \bar{K}_{n_i} where \bar{K}_{n_i} is the complement of the complete graph K_{n_i} .

Proof: Denote the vertex set of $EDG(T_m)$ as $V =$

$\{v_1, v_2, \dots, v_{3m+2}, v_1', v_2', \dots, v_{3m+2}'\}$

Using algorithm, the colors in $EDG(T_m)$ are classified as follows :

The color class

$$C_1 = \begin{cases} v_1 \cup v_1' & \text{if } m = 1 \\ v_1 \cup v_1' \cup (\cup_{i=2}^n v_{3i} \cup v_{3i}') & \text{if } 2 \leq n \leq m \end{cases}$$

The color class

$$C_2 = \begin{cases} v_2 & \text{if } m = 1 \\ v_2 \cup v_7' & \text{if } m = 2 \\ v_2 \cup v_7' \cup (\cup_{i=3}^n v_{3i+1} \cup v_{3i+1}') & \text{if } 3 \leq n \leq m \end{cases}$$

The color class

$$C_3 = \begin{cases} v_2' & \text{if } m = 1 \\ v_2' \cup v_7 & \text{if } m = 2 \\ v_2' \cup v_7 \cup (\cup_{i=3k/k \in \mathbb{N}}^n v_{3i+2} \cup v_{3i+2}') & \text{if } 3 \leq n \leq m \end{cases}$$

The color class $C_4 = \begin{cases} v_3 \cup v_3' & \text{if } m = 1 \\ v_3 \cup v_3' \cup (\cup_{i=3k-1/k \in \mathbb{N}}^n v_{3i+2} \cup v_{3i+2}') & \text{if } 2 \leq n \leq m \end{cases}$

The color class $C_5 = \begin{cases} v_4 \cup v_4' & \text{if } m \leq 3 \\ v_4 \cup v_4' \cup (\cup_{i=3k+1/k \in \mathbb{N}}^n v_{3i+2} \cup v_{3i+2}') & \text{if } 4 \leq n \leq m \end{cases}$

The color class $C_6 = v_5 \cup v_5' \quad \forall m$.

Let n_i be the number of colors in C_i where $1 \leq i \leq 6$ respectively in the $EDG(T_m)$ graph.

Thus the vertices in the induced subgraph obtained from each of the color classes are not adjacent, since they are not adjacent in $EDG(T_m)$.

Hence the induced subgraph of n_i vertices relative to each of the above color class is a totally disconnected graph \bar{K}_{n_i} , where \bar{K}_{n_i} is the complement of the complete graph K_{n_i} , $1 \leq i \leq 6$.

Theorem 3.3. The conflict-free chromatic number of $\text{EDG}(T_m)$ is $\Delta + 1$.

Proof: Consider the $\text{EDG}(T_m)$ which has $6m+4$ vertices. Denote them as

$$V = \{v_1, v_2, \dots, v_{3m+2}, v'_1, v'_2, \dots, v'_{3m+2}\} = V_1 \cup V_2$$

$$\text{where } V_1 = \{v_1, v_2, v_3, v_4, v_5, v'_1, v'_2, v'_3, v'_4, v'_5\} \text{ and } V_2 = V \setminus V_1.$$

Case (i) : $m = 1$

Consider the vertex set V_1 of $\text{EDG}(T_1)$. Since $\deg(v_2) = \deg(v'_2), |N(v_2)| = |N(v'_2)| = \Delta$ and $|N(v_i)| = |N(v'_i)| < \Delta$ where $1 \leq i \leq 5, i \neq 2$. Hence $\Delta + 1$ colors are needed to color $N[v_2]$.

Since $N(v_i) \cap N(v'_i) = \emptyset$, $N[v_2]$ and $N[v'_2]$ receives same colors whereas v_2 and v'_2 receives different colors among the $\Delta + 1$ colors.

Hence, $\text{EDG}(T_1)$ is $\Delta + 1$ conflict-free colorable.

Case (ii) : $m \geq 2$ and m is finite

Using case (i) the vertices of V_1 are colored. The vertices in V_2 are colored as follows.

For $n = 2, 3, 4, \dots, m$, let $l_n = \{3n-4, 3n, 3n+1, 3n+2\}$ and $S_n = \{3n-1\}$. Now

$$N(v_i) \cap N(v_j) = \begin{cases} \{v_k\} & \text{if } i, j \in l_n, i \neq j \text{ and } k \in S_n \\ \phi & \text{otherwise} \end{cases}$$

$$N(v'_i) \cap N(v'_j) = \begin{cases} \{v_k\} & \text{if } i, j \in l_n, i \neq j \text{ and } k \in S_n \\ \phi & \text{otherwise} \end{cases}$$

Also $N(v_s) \cap N(v'_t) = \emptyset$ for $1 \leq s, t \leq 3m+2$.

Thus the colors used to color these vertices in V_1 is enough to color the vertices of V_2 .

Hence $\text{EDG}(T_m)$ is $\Delta + 1$ conflict-free colorable.

4. Conclusion

In this paper, we proved the existence of conflict-free coloring in $\text{EDG}(T_m)$ and the induced subgraph of each of the color class is a totally disconnected graph. Also we proved that the conflict-free chromatic number of $\text{EDG}(T_m)$ is $\Delta + 1$.

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